

Universal Cycles for Weak Orders

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Abstract

Universal cycles are generalizations of de Bruijn cycles and Gray codes that were introduced originally by Chung, Diaconis, and Graham in 1992. They have been developed by many authors since, for various combinatorial objects such as strings, subsets, permutations, partitions, vector spaces, and designs. One generalization of universal cycles, which require almost complete overlap of consecutive words, is s -overlap cycles, which relax such a constraint. In this paper we study weak orders, which are relations that are transitive and complete. We prove the existence of universal and s -overlap cycles for weak orders, as well as for fixed height and/or weight weak orders, and apply the results to cycles for ordered partitions as well.

1 Introduction

A **weak order on** $[n]$ is a relation \preceq that is transitive and complete. We write $x \equiv y$ if $x \preceq y$ and $y \preceq x$, and we write $x \prec y$ if $x \preceq y$ but $y \not\preceq x$. A weak order on $[n]$ can be written as a permutation of $[n]$ with consecutive symbols separated by \equiv or \prec . (See [13], fasc. 2, problem 105.) We use the notation $\mathcal{W}(n)$ to represent the set of all weak orders on $[n]$. For example, $\mathcal{W}(3)$ contains the weak orders:

$$\begin{aligned} 1 \equiv 2 \equiv 3, & \quad 1 \equiv 2 \prec 3, \quad 1 \prec 2 \equiv 3, \quad 1 \equiv 3 \prec 2, \\ 2 \equiv 3 \prec 1, & \quad 2 \prec 1 \equiv 3, \quad 3 \prec 1 \equiv 2, \quad 1 \prec 3 \prec 2, \\ 2 \prec 1 \prec 3, & \quad 1 \prec 2 \prec 3, \quad 3 \prec 1 \prec 2, \quad 2 \prec 3 \prec 1, \\ 3 \prec 2 \prec 1. & \end{aligned}$$

We are primarily interested in finding universal cycles for $\mathcal{W}(n)$ and some its subsets. Let \mathcal{C} be a set of k strings, each of length n . A **universal cycle** (ucycle) $a_0 a_1 \dots a_{k-1}$ for \mathcal{C}

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is a word such that each object $c \in \mathcal{C}$ appears exactly once as a subword $a_{i+1}a_{i+2} \dots a_{i+n}$, where subscripts are taken modulo k . That is, we require objects at the end of the string to wrap-around to the beginning; for example some object will appear as the subword $a_{k-1}a_ka_0a_1 \dots a_{n-3}$. Note in particular that every $a_{i+1}a_{i+2} \dots a_{i+n}$ is a string in \mathcal{C} . For example, a ucycle for the set of binary strings of length three is the de Bruijn cycle

$$00010111.$$

Ucycles are generalizations of de Bruijn cycles and Gray codes that were originally introduced by Chung, Diaconis, and Graham in 1992 [2]. They have been developed for various combinatorial objects, such as binary strings, subsets, restricted multisets, permutations, partitions, lattice paths, and designs, by many authors since (see [1, 2, 4, 7, 8, 9, 10, 11]). Note that this definition requires our subwords to be contiguous. Others have considered variations in which subwords need not be contiguous, such as rosaries [6], or combs [3].

Some combinatorial objects do not readily lend themselves to the universal cycle structure. For example, all permutations of $[n]$ written as strings can never be listed in a universal cycle because the first $n - 1$ letters will always completely determine the last letter. This forces a set of multiple disjoint cycles instead of one long cycle. To deal with this problem, we may wish to either use a different representation for our set of objects, or generalize the concept of a ucycle. Returning to our permutation example, we can define an alternate representation using the first $n - 1$ letters of each permutation since the last letter is completely determined by the first $n - 1$. We can generalize ucycles by using overlap cycles, first defined in [5]. An **s -overlap cycle** (ocycle) is an ordered listing of the objects so that the last s letters of one word are the first s letters of its successor in the listing. Note that an $(n - 1)$ -ocycle is a ucycle.

In this paper we prove that ucycles exist for all weak orders on $[n]$ (Theorem 3.2), as well as for those of fixed weight and/or fixed height (Theorems 3.3, 3.7, and 3.8), and for those over a fixed multiset (Lemma 3.4), where the relevant terms are defined below. We also apply these results to obtain ucycles for various types of ordered partitions (Theorems 5.2 and 5.3), and construct s -ocycles for various sets of weak orders as well (Theorem 4.1, Theorem 4.3, Theorem 4.5, Theorem 4.6).

2 Definitions

For each word $w \in \mathcal{W}(n)$ we can define the **height** of an element $j \in [n]$ in the word w to be the number of symbols \prec that precede it in the weak order. This gives us an alternative representation for $w \in \mathcal{W}(n)$ by a word $w_1w_2 \dots w_n$ where letter w_j is the height of element

j in the word w . The set $\mathcal{W}(3)$ listed above contains the corresponding words:

$$\begin{array}{cccc} 000, & 001, & 011, & 010, \\ 100, & 101, & 110, & 021, \\ 102, & 012, & 120, & 201, \\ 210. \end{array}$$

We will utilize this word representation of each weak order, and so we write $w = w_1 w_2 \dots w_n$. Note that a weak order on $[n]$ in this representation will always have length n .

When representing weak orders as words, we will use various types of abbreviations. We define an exponential notation to represent repeated elements, i.e.

$$a^i = \overbrace{a \dots a}^{i \text{ times}}.$$

Also, we will represent runs of consecutive elements as:

$$[i, j] = i(i+1)(i+2) \dots (j-1)j.$$

We will sometimes refer to the **height** of a weak order $w \in \mathcal{W}(n)$. This is defined to be $\text{ht}(w) = \max\{h \mid h \text{ is the height of } j \in [n] \text{ in } w\}$. In our word representation, height corresponds to the maximum letter in the word $w_1 w_2 \dots w_n$. Note that this is equivalent to defining $\{w_1, w_2, \dots, w_n\} = \{0, 1, 2, \dots, \text{ht}(w)\}$, with multiplicities removed. It is often useful to restrict our attention to the subset of $\mathcal{W}(n)$ with a specific or fixed height. Let $0 \leq h < n$, and define

$$\mathcal{W}(n, h) = \{w \in \mathcal{W}(n) \mid \text{ht}(w) = h\}.$$

We will also consider the subsets of $\mathcal{W}(n)$ that correspond to a fixed multiset or fixed weight. Define the **multiset** of a weak order $w = w_1 w_2 \dots w_n$, or $\text{ms}(w)$, to be the unordered multiset of elements $\{w_1, w_2, \dots, w_n\}$. For a fixed multiset M , we define

$$\mathcal{W}_M(n) = \{w \in \mathcal{W}(n) \mid \text{ms}(w) = M\}.$$

The **weight** of a weak order $w = w_1 w_2 \dots w_n$ is the sum of its letters. We denote this by

$$\text{wt}(w) = \sum_{i=1}^n w_i.$$

To identify the subset of $\mathcal{W}(n)$ that contains only weak orders with weight k , we write $\mathcal{W}_k(n)$. When considering $w = w_1 w_2 \dots w_n \in \mathcal{W}_k(n)$, note that $w_n = k - \sum_{i=1}^{n-1} w_i$, and so the first $n-1$ letters of w completely define it. We define the **prefix** of a word w to be $w^- = w_1 w_2 \dots w_{n-1}$. When w^- completely identifies w as in words from $\mathcal{W}_k(n)$, this is called the **prefix representation**. Similarly, w^+ , or the **suffix** of w , is defined to be $w_2 w_3 \dots w_n$.

3 Ucycle Results

Most of our results are obtained by finding Euler tours in graphs. An Euler tour is a closed circuit that contains each edge of the graph exactly once. We will use the following theorem often.

Theorem 3.1. ([14], p. 60) *A directed graph G is eulerian if and only if it is both balanced and weakly connected.*

We begin with the most general result.

Theorem 3.2. *For all $n \in \mathbb{Z}^+$, there exists a ucycle for $\mathcal{W}(n)$.*

Proof. Fix $n \in \mathbb{Z}^+$ and define the graph $G(n)$ as follows:

$$V(G(n)) = \{v = v_1 \dots v_{n-1} \mid v = w^- \text{ for some } w \in \mathcal{W}(n)\}$$

and

$$E(G(n)) = \{(v_1, v_2) \mid v_1 = w^- \text{ and } v_2 = w^+ \text{ for some } w \in \mathcal{W}(n)\}.$$

Note that the edge set is well-defined, since any prefix w^- of some $w \in \mathcal{W}(n)$ is also the suffix of the word $w_n w_1 w_2 \dots w_{n-1} \in \mathcal{W}(n)$. We will explicitly show a method to construct a ucycle for $\mathcal{W}(n)$ using the graph $G(n)$ for $n \geq 2$. Note that for $n = 1$, the ucycle is the single letter 0.

By the construction of $G(n)$, it is clear that an Euler tour will correspond to a ucycle. Using Theorem 3.1, we need only show that the graph is weakly connected and balanced, and we are done. The graph is clearly balanced, since any incoming edge $(w_1 \dots w_{n-1}, w_2 \dots w_n)$ can be paired with the outgoing edge $(w_2 \dots w_n, w_3 \dots w_n w_1)$ at the vertex $w_2 \dots w_n$. Next, we will show a path from any vertex $w_1 \dots w_{n-1}$ to the vertex $00 \dots 0$. We may apply the rotation function ρ , which maps $w_1 w_2 \dots w_n$ to $w_2 \dots w_n w_1$, as many times as necessary until we arrive at some $u = u_1 u_2 \dots u_n \in \mathcal{W}(n)$ with u_n equal to the height of w . Then u is represented in $G(n)$ as the edge $(u_1 \dots u_{n-1}, u_2 \dots u_n)$. Since u_n is a maximum letter in w , we must also have $u_1 \dots u_{n-1} 0 \in \mathcal{W}(n)$, and so we have $(u_1 \dots u_{n-1}, u_2 \dots u_{n-1} 0) \in E(G(n))$. Note that $u_2 \dots u_{n-1} 0$ has more zeros than $w_1 \dots w_{n-1}$. Repeating this process, we add more zeros at every step, which eventually must terminate when we arrive at the vertex with $n - 1$ zeros. Thus, we have a path from $w_1 \dots w_{n-1}$ to $00 \dots 0$, and so G is weakly connected. □

Next, we consider some subsets of $\mathcal{W}(n)$. Define

$$\mathcal{W}_k^-(n) = \{w^- \mid w \in \mathcal{W}_k(n)\}.$$

Theorem 3.3. *For all $n, k \in \mathbb{Z}^+$ with $k \leq \binom{n}{2}$, there exists a ucycle for $\mathcal{W}_k^-(n)$.*

To prove this theorem, we use the same approach as in the previous proof but need a few lemmas first to simplify our method.

Lemma 3.4. *Let $n \in \mathbb{Z}^+$, and let M be some fixed multiset. Define the set A to be the set of all permutations of M . Then there exists a ucycle for A using the prefix representation.*

Proof. Construct a graph G_M with

$$V(G_M) = \{v = a_1 a_2 \dots a_{n-2} \mid v = a^- \text{ for some } a \in A^-\}$$

and

$$E(G_M) = \{(v_1, v_2) \mid v_1 = a^- \text{ and } v_2 = a^+ \text{ for some } a \in A^-\}.$$

Note that the edges in this graph correspond to the elements of A^- , so we would like to find an Euler tour in G_M , which will produce a ucycle as desired.

For any vertex $w_1 \dots w_{n-4} w_{n-3} w_{n-2}$ corresponding to word $w = w_1 \dots w_n$, there is an undirected path to the vertex $w_1 \dots w_{n-4} w_{n-2} w_{n-3}$. This path is:

$$\begin{aligned} w_1 \dots w_{n-4} w_{n-3} w_{n-2} &\leftarrow w_n w_1 \dots w_{n-4} w_{n-3} \\ &\leftarrow w_{n-1} w_n w_1 \dots w_{n-4} \\ &\rightarrow w_n w_1 \dots w_{n-4} w_{n-2} \\ &\rightarrow w_1 \dots w_{n-4} w_{n-2} w_{n-3}. \end{aligned}$$

Thus we can always find an undirected path from one vertex to another whose difference is the transposition of elements $n-1$ and $n-2$. Since we also have paths from a vertex $v = w^{--}$ to $\rho(v) = \rho(w)^{--}$ (where ρ is the previously defined rotation function), we can find an undirected path between any two vertices that differ by adjacent transpositions. Then, since the set of all adjacent transpositions generate all permutations of a set (see [12]), there exists an undirected path between any vertices with the same multiset, so the graph G_M is weakly connected.

Note that the graph G_M must also be balanced, since having fixed multiset ensures that any edge

$$w_1 \dots w_{n-2} \rightarrow w_2 \dots w_{n-1}$$

can be balanced by the edge

$$w_{n-1} w_1 \dots w_{n-3} \rightarrow w_1 \dots w_{n-2}.$$

Thus the graph is balanced and connected, and so is eulerian by Theorem 3.1.

□

To prove Theorem 3.3, we will construct a new transition graph, $G_k(n)$. Define

$$V(G_k(n)) = \{v = w_1 \dots w_{n-2} \mid v = w^- \text{ for some } w \in \mathcal{W}_k^-(n)\}$$

and

$$E(G_k(n)) = \{(v_1, v_2) \mid v_1 = w^- \text{ and } v_2 = w^+ \text{ for some } w \in \mathcal{W}_k^-(n)\}.$$

As before, note that this definition of edges is well-defined, for if $w_1 w_2 \dots w_{n-1} \in \mathcal{W}_k^-(n)$, then $w_{n-1} w_1 w_2 \dots w_{n-2} \in \mathcal{W}_k^-(n)$.

Since the vertices in $G_k(n)$ are words of length $n - 2$, we identify the vertex v that is the minimum word in lexicographic order as the **minimum vertex**. It will be useful to determine exactly what word v is, and so we have the following fact, together with a lemma.

Fact 3.5. ([13], fasc. 3, p. 19) *For every $k \in \mathbb{Z}^+$, there are unique $a, b \in \mathbb{Z}^+$ with $a > b$ so that $k = \binom{a}{2} + \binom{b}{1}$.*

Lemma 3.6. *Fix $n, k \in \mathbb{Z}^+$. Define*

$$w = [0, b - 1]b^2[b + 1, a]$$

with $k = \binom{a}{2} + \binom{b}{1}$, and adding additional 0's to the front to create a word of length n . Then the minimum vertex v in $G_k(n)$ is $v = w^{--}$.

Proof. If we are looking for the element $w \in \mathcal{W}_k(n)$ with w^{--} minimum in lex order, we may consider only the elements of $\mathcal{W}_k(n)$ with the largest two elements in positions $n - 1$ and n . These elements must be as large as possible in order to get the smallest elements possible in positions 1 to $n - 2$, so the string we desire will have letters of the largest height possible. Note that if $\text{ht}(w) = h$ and $w \in \mathcal{W}_k(n)$, then we must have $k \geq \binom{h+1}{2}$. Thus we choose w_n as large as possible so that $\binom{w_n}{2} \leq k < \binom{w_n+1}{2} = \binom{w_n}{2} + w_n$.

Now, by our choice of w_n , the remaining weight is $0 \leq k - \binom{w_n}{2} < w_n$. Thus $k - \binom{w_n}{2} \in \{0, 1, \dots, w_n\}$, and so $k - \binom{w_n}{2}$ is equal to some element $b \leq m$. Then we can add another letter b to the word to obtain a word with multiset of at most n elements. If the cardinality of the multiset is less than n , we add 0's to reach cardinality n . Sorting this multiset from smallest element to largest we obtain the desired weak order w . \square

Before we begin the proof, we define one more term. We say that there is a **duplicate at index i** if $w_i = w_{i+1} > 0$ for some $w_1 w_2 \dots w_n \in \mathcal{W}(n)$. Now we are finally ready to prove Theorem 3.3.

Proof of Theorem 3.3. First, note that if $k > \binom{n}{2}$, then it is not possible to construct a weak order on $[n]$ with weight k . The maximum weight weak order possible is one with multiset $\{0, 1, \dots, n - 1\}$, which must have weight

$$\sum_{i=0}^{n-1} i = \binom{n}{2}.$$

We would like to show an undirected path from any vertex $x^{--} = x_1x_2 \dots x_{n-2}$ to the minimum vertex $v^{--} = v_1v_2 \dots v_{n-2}$ in $G_k(n)$. Since $x \in \mathcal{W}_k(n)^{--}$, we may define x_{n-1}, x_n accordingly so that $x \in \mathcal{W}_k(n)$. By Lemma 3.4, we may assume that $x_1 \leq x_2 \leq \dots \leq x_{n-2} \leq x_{n-1} \leq x_n$. Our first observation is that if x has either no duplicates or one duplicate then we must have $x = v$, and hence $x^{--} = v^{--}$. This follows from the fact that the minimum vertex writes weight k in the most compact way possible, i.e. with either zero or one duplicate. Now we assume that x has at least two duplicates, and we have several cases depending on the relationship between $\text{ht}(v)$ and $\text{ht}(x)$.

If $\text{ht}(x) = \text{ht}(v) = h$, then using Lemma 3.4 we may rewrite both weak orders as

$$x = [0, h]x_{h+2}x_{h+3} \dots x_n \text{ and } v = [0, h]v_{h+2}v_{h+3} \dots v_n$$

where $x_{h+2} \leq x_{h+3} \leq \dots \leq x_n$ and $v_{h+2} \leq v_{h+3} \leq \dots \leq v_n$. Now since $x \neq v$ there must be indices $i < j$ with $i, j \in \{h+2, h+3, \dots, n\}$ so that $x_i > v_i$ and $x_j < v_j$. Using Lemma 3.4 again, we can reorder the letters of x^{--} so that we have

$$x^{--} = x_i x_j [0, h]x_{h+2}x_{h+3} \dots x_{n-2}.$$

Then we have the undirected path:

$$\begin{aligned} x_i x_j [0, h]x_{h+2}x_{h+3} \dots x_{n-2} &\rightarrow x_j [0, h]x_{h+2}x_{h+3} \dots x_{n-2}x_{n-1} \\ &\rightarrow [0, h]x_{h+2}x_{h+3} \dots x_n \text{ (with } x_i, x_j \text{ missing)} \\ &\leftarrow (x_j + 1)[0, h]x_{h+2}x_{h+3} \dots x_{n-1} \\ &\leftarrow (x_i - 1)(x_j + 1)[0, h]x_{h+2}x_{h+3} \dots x_{n-2} \end{aligned}$$

Continuing in this manner, we will eventually arrive at the vertex v^{--} .

If $\text{ht}(x) < \text{ht}(v)$, then we consider the fact that x must have at least two duplicates, say at x_i and x_j . Using Lemma 3.4, rewrite x^{--} with x_i, x_j at the front, i.e.

$$x^{--} = x_i x_j x_1 x_2 \dots x_{n-2}.$$

Then we have the undirected path:

$$\begin{aligned} x_i x_j x_1 x_2 \dots x_{n-2} &\rightarrow x_j x_1 x_2 \dots x_{n-2}x_{n-1} \\ &\rightarrow x_1 x_2 \dots x_{n-2}x_{n-1}x_n \\ &\leftarrow (x_j + 1)x_1 x_2 \dots x_{n-2}x_{n-1} \\ &\leftarrow (x_i - 1)(x_j + 1)x_1 x_2 \dots x_{n-2} \end{aligned}$$

We continue to decrease x_i and increase x_j until either $x_i = 0$ or $x_j = \text{ht}(x) + 1$. In either case, we have removed a duplicate. Continuing, we will eventually arrive at a vertex with either 0 or 1 duplicates, which as stated previously must be the minimum vertex v^{--} .

Next, we note that it is not possible to have $\text{ht}(x) > \text{ht}(v)$, as otherwise x would have been chosen as the minimum vertex (v was chosen so as to have maximum height). Thus

in all cases we have shown an undirected path from x^{--} to v^{--} , so the graph must be weakly connected.

Lastly, the graph must be balanced, since any outgoing edge from vertex $w_1 \dots w_{n-2}$ to $w_2 \dots w_{n-1}$ can be paired with the incoming edge $w_{n-1}w_1 \dots w_{n-3}$ to $w_1 \dots w_{n-2}$ and this pairing gives a unique incoming edge for each outgoing edge. \square

Next, we show that there is always a universal cycle for $\mathcal{W}(n, h)$, the set of weak orders on $[n]$ with height h . We will restrict our attention to the case when $h < n - 1$, for if $h = n - 1$ then we are considering the set of all permutations of an n -set. Ucycles for permutations can be found easily using the prefix representation, or for alternative methods see [8, 11].

Theorem 3.7. *For all $n \in \mathbb{Z}^+$ and all $h \in \mathbb{N}$ with $0 \leq h < n - 1$, there exists a universal cycle for $\mathcal{W}(n, h)$.*

Proof. We construct the standard transition graph $G(n, h)$ as follows. We define

$$V(G(n, h)) = \mathcal{W}^-(n, h) = \mathcal{W}^+(n, h)$$

and

$$\begin{aligned} E(G(n, h)) &= \{(v, w) \mid v \in \mathcal{W}^-(n, h), w \in \mathcal{W}^+(n, h), \text{ and } v_i = w_{i+1} \text{ for } 1 \leq i < n - 1\} \\ &= \mathcal{W}(n, h). \end{aligned}$$

Note that if $w_1 w_2 \dots w_n \in \mathcal{W}(n, h)$ then $w_n w_1 w_2 \dots w_{n-1} \in \mathcal{W}(n, h)$ and so $G(n, h)$ is balanced.

To finish the proof, we must show that the graph is connected. Define our minimum vertex in the graph to be $v^- = 0^{n-h-1}[1, h]$. Let $x^- = x_1 x_2 \dots x_{n-1}$ be an arbitrary vertex, and define x_n so that $x = x_1 x_2 \dots x_{n-1} x_n \in \mathcal{W}(n, h)$. We will show a path from x^- to the minimum vertex v^- by first illustrating that all permutations of the minimum vertex are connected, and then describing a path from x^- to some permutation of v^- .

Starting from the vertex $v^- = 0^{n-h-1}[1, h]$, we show that any adjacent transposition can be reached by a path. Since adjacent transpositions generate all permutations, this proves that all permutations of v^- are connected. Let $w_1 w_2 \dots w_{n-1}$ be a permutation of v^- . We will show that the letters w_i and w_{i+1} may be transposed. First, we note that $\{w_1, w_2, \dots, w_{n-1}\} = \{0, 1, 2, \dots, h\}$ by the definition of v^- . Then we find the path as follows:

$$\begin{aligned} w_1 w_2 \dots w_{n-1} &\rightarrow w_2 w_3 \dots w_{n-1} w_i \\ &\rightarrow w_3 w_4 \dots w_{n-1} w_i w_1 \\ &\vdots \\ &\quad (\text{rotations of the weak order } w_1 w_2 \dots w_{n-1} w_i) \end{aligned}$$

$$\begin{aligned}
&\rightarrow w_{i+1}w_{i+2} \dots w_{n-1}w_iw_1w_2 \dots w_{i-1} \\
&\rightarrow w_{i+2}w_{i+3} \dots w_{n-1}w_iw_1w_2 \dots w_{i-1}w_{i+1} \\
&\rightarrow w_{i+3}w_{i+4} \dots w_{n-1}w_iw_1w_2 \dots w_{i-1}w_{i+1}w_i \\
&\vdots \quad (\text{rotations of the weak order } w_1w_2 \dots w_{i-1}w_{i+1}w_iw_{i+2}w_{i+3} \dots w_{n-1}w_i) \\
&\rightarrow w_1w_2 \dots w_{i-1}w_{i+1}w_iw_{i+2}w_{i+3} \dots w_{n-1}
\end{aligned}$$

Note that along this path, every edge contains all letters from the set $\{w_1, w_2, \dots, w_{n-1}\} = \{0, 1, 2, \dots, h\}$, and so is a valid weak order from $\mathcal{W}(n, h)$.

Now let $x^- = x_1x_2 \dots x_{n-1} \in V(G(n, h))$ be arbitrary. We want to define a path from x^- to some permutation of the minimum vertex in $G(n, h)$ by repeatedly replacing any duplicates in x with 0. To create this path, we first define x_n so that $x = x_1x_2 \dots x_n \in \mathcal{W}(n, h)$. Then if x has a duplicate at index i , we can replace it by following the path:

$$\begin{aligned}
x_1x_2 \dots x_{n-1} &\rightarrow x_2x_3 \dots x_{n-1}x_n \\
&\vdots \quad (\text{rotations of } x) \\
&\rightarrow x_ix_{i+1} \dots x_nx_1x_2 \dots x_{i-2} \\
&\rightarrow x_{i+1}x_{i+2} \dots x_nx_1x_2 \dots x_{i-1} \\
&\rightarrow x_{i+2}x_{i+3} \dots x_nx_1x_2 \dots x_{i-1}0
\end{aligned}$$

Repeating this procedure, we will eventually arrive at a vertex that is the prefix of some weak order y that is a permutation of $0^{n-h}[1, h]$. Following rotations, we will eventually arrive at a vertex that is a permutation of v^- . Thus there exists a path from x^- to v^- .

By Theorem 3.1, since $G(n, h)$ is balanced and connected, it is eulerian. Therefore we can find a ucycle by following the Euler tour in $G(n, h)$. \square

Finally, we prove the following result on the subset of fixed weight, fixed height weak orders on $[n]$.

Theorem 3.8. *For every $n, k, h \in \mathbb{Z}^+$ with $k \leq \binom{n}{2}$ and $0 \leq h < n$, there is a ucycle for $\mathcal{W}_k^-(n, h)$.*

Proof. We construct the transition graph $G_k(n, h)$ as usual, with

$$V(G_k(n, h)) = \mathcal{W}_k^{--}(n, h)$$

and

$$E(G_k(n, h)) = \mathcal{W}_k^-(n, h).$$

First, we note that the graph is even, since if $w_1w_2 \dots w_{n-1} \in \mathcal{W}_k^-(n, h)$, then $w_2 \dots w_{n-1}w_1 \in \mathcal{W}_k^-(n, h)$.

Next we must show that the graph is connected. We define the minimum vertex $\mathbf{v} = v_1v_2 \dots v_{n-2}$ by constructing a specific weak order \mathbf{w} in $\mathcal{W}_k(n, h)$, and then removing the

last two elements. Any weak order in $\mathcal{W}_k(n, h)$ must contain the letters $0, 1, 2, \dots, h$, which have total weight $k' = \sum_{i=0}^h i$. Let $\mathbf{s} = s_1 s_2 \dots s_{n-(h+1)}$ be the lexicographically minimum element of the set $\mathcal{B}_{k-k'}^{h+1}(n - (h+1))$, the set of words of length $n - (h+1)$ and weight $k - k'$ using the alphabet $\{0, 1, 2, \dots, h\}$. Then extend s to be a weak order by defining $\mathbf{w} = s_1 s_2 \dots s_{n-(h+1)} 0 1 \dots h$. At this point we have $\mathbf{w} \in \mathcal{W}_k^-(n, h)$, but we reorder the letters of \mathbf{w} so that

$$\mathbf{w} = [0, h] w_{h+2} w_{h+3} \dots w_n \text{ where } w_{h+2} \leq w_{h+3} \leq \dots \leq w_n.$$

Then define the minimum vertex to be $\mathbf{v} = w_1 w_2 \dots w_{n-2}$.

Now we consider an arbitrary vertex $\mathbf{x} = x_1 x_2 \dots x_{n-2}$, and we will show that it is connected to the minimum vertex. Since $\mathbf{x} \in \mathcal{W}_k^-(n, h)$, we define x_{n-1} and x_n so that $x_1 x_2 \dots x_n \in \mathcal{W}_k(n, h)$. Then, by Lemma 3.4, we know that \mathbf{x} must be connected to some vertex

$$\mathbf{y} = [0, h] y_{h+2} y_{h+3} \dots y_{n-2} \text{ where } y_{h+2} \leq y_{h+3} \leq \dots \leq y_{n-2}.$$

If $\mathbf{y} = \mathbf{v}$, then we are done. Otherwise, there exists some $i, j \in \{h+2, h+3, \dots, n-2\}$ so that $y_i > v_i$ and $y_j < v_j$. Using rotations and Lemma 3.4, there exists a path in $G_k(n, h)$ from \mathbf{y} to the vertex

$$y_i y_j y_1 y_2 \dots y_{i-1} y_{i+1} \dots y_{j-1} y_{j+1} \dots y_{n-2}.$$

Define y_{n-1} and y_n so that $y_1 y_2 \dots y_n \in \mathcal{W}_k(n, h)$. Then we construct the following path in $G_k(n, h)$:

$$\begin{aligned} & y_i y_j y_1 y_2 \dots y_{i-1} y_{i+1} \dots y_{j-1} y_{j+1} \dots y_{n-2} \\ \rightarrow & y_j y_1 y_2 \dots y_{i-1} y_{i+1} \dots y_{j-1} y_{j+1} \dots y_{n-2} y_{n-1} \\ \rightarrow & y_1 y_2 \dots y_{i-1} y_{i+1} \dots y_{j-1} y_{j+1} \dots y_{n-1} y_n \\ \rightarrow & y_2 \dots y_{i-1} y_{i+1} \dots y_{j-1} y_{j+1} \dots y_n v_i \\ \rightarrow & y_3 \dots y_{i-1} y_{i+1} \dots y_{j-1} y_{j+1} \dots y_n v_i v_j \end{aligned}$$

Reordering (by Lemma 3.4) and shuffling/replacing elements, we can find a path to the vertex

$$y_1 y_2 \dots y_{i-1} v_i y_{i+1} \dots y_{j-1} v_j y_{j+1} \dots y_{n-2}.$$

Note that by requiring that both \mathbf{v} and \mathbf{y} start with $[0, h]$, we are ensured that these intermediary vertices always contain the set $\{0, 1, 2, \dots, h\}$, and hence all edges represent valid weak orders in $\mathcal{W}_k^-(n, h)$. Continuing this process, eventually we will arrive at \mathbf{v} . Since the graph is even and connected, it must be eulerian by Theorem 3.1. \square

4 Overlap Cycle Results

Many of our ucycle results for weak orders have corresponding ocycle results.

Theorem 4.1. *For all $n \in \mathbb{Z}^+$ and for all $s \in \mathbb{Z}^+$ with $1 \leq s \leq n-1$, there is an s -ocycle for $\mathcal{W}(n)$.*

To prove this, we have two cases: $1 \leq s \leq \frac{n}{2}$ (vertices do not overlap to make a weak order), and $\frac{n}{2} < s \leq n-1$ (vertices must overlap to make a weak order). First, we define the s -prefix, w^{s-} , and s -suffix, w^{s+} , of a word as:

$$w^{s-} = w_1 w_2 \dots w_s$$

and

$$w^{s+} = w_{n-s+1} w_{n-s+2} \dots w_n.$$

Fact 4.2. *If $w = w_1 w_2 \dots w_n \in \mathcal{W}(n)$, then any permutation w' of the letters of w is also in $\mathcal{W}(n)$.*

In light of this fact, we note that $\{w^{s-} \mid w \in \mathcal{W}(n)\} = \{w^{s+} \mid w \in \mathcal{W}(n)\}$. In fact, these two sets are also equal to the set

$$\mathcal{W}^s(n) = \{w = w_1 w_2 \dots w_s \mid w \text{ is a subword of some } w' \in \mathcal{W}(n)\}.$$

Proof of Theorem 4.1. We define a transition graph $G^s(n)$ as follows. Let $V(G^s(n))$ equal to the set of all possible overlaps, i.e. $V(G^s(n)) = \mathcal{W}^s(n)$. Define $E(G^s(n))$ to contain one edge for each weak order by creating a directed edge (u, v) for each weak order w that begins with u and ends with v . We will show that this graph contains an Euler tour, which will give us an s -ocycle for $\mathcal{W}(n)$.

First, $G^s(n)$ must have $d^+(u) = d^-(u)$ for all $u \in V(G^s(n))$, since any permutation of a weak order is again a weak order. That is, if $u = u_1 u_2 \dots u_s$ is a prefix of some weak order $u_1 u_2 \dots u_n$, then we have incoming edge $u_{s+1} u_{s+2} \dots u_n u_1 u_2 \dots u_s$ and outgoing edge $u_1 u_2 \dots u_s u_{s+1} \dots u_n$. Thus we can pair together each incoming edge with an outgoing edge, so we must have $d^+(u) = d^-(u)$.

Lastly we must show that $G^s(n)$ is connected. We will show that any vertex v is connected to the vertex 0^s . Let $v = v_1 v_2 \dots v_s \in V(G^s(n))$ with $h = \text{ht}(v)$. If $s \leq \frac{n}{2}$, then we must have an edge (v, u) to some vertex u with $\text{ht}(u) < h$. If $s > \frac{n}{2}$, the weak order $v_1 v_2 \dots v_s v_{s+1} \dots v_n$ is represented by the edge $(v, v_{n-s} v_{n-s+1} \dots v_n)$. Note that $\text{ht}(v_{n-s} v_{n-s+1} \dots v_n) \leq h$, and that any letter v_i of maximum height in $v_{n-s} v_{n-s+1} \dots v_n$ must have $n-s \leq i \leq s$. Thus by repeating this procedure (at most n times if $s = n-1$ and v_{n-1} has maximum height in v), we reach some vertex $u \in V(G^s(n))$ with $\text{ht}(u) < h$.

In either case, we have moved to a vertex with smaller height. By repeating this procedure, we will eventually reach a vertex with height 0. The only vertex with height 0 is the vertex 0^s , and so we have arrived at our destination.

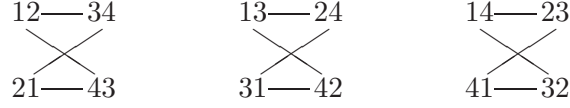


Figure 1: Transition Graph: 2-ocycles for permutations of $\{1, 2, 3, 4\}$

Since our graph $G^s(n)$ is connected and $d^+(v) = d^-(v)$ for all $v \in V(G^s(n))$, we must have an Euler tour in the graph. This Euler tour will translate to an s -ocycle on $\mathcal{W}(n)$. \square

When we consider s -ocycles for fixed weight weak orders on $[n]$, we notice that the following theorems follow immediately from their corresponding results about ucycles with very small adjustments. Note, however, that we must consider whether or not s and n are relatively prime. For example, if we consider all permutations of the set $\{1, 2, 3, 4\}$, the transition graph for 2-ocycles is disconnected, as shown in Figure 1.

Lemma 4.3. *Let $n, s \in \mathbb{Z}^+$ with $n \geq 2$ and $1 \leq s \leq n-2$ and let M be some fixed multiset. Define the set A to be the set of all permutations of M .*

1. *If $\gcd(s, n) = 1$, then there is an s -ocycle for A .*
2. *If $\gcd(s, n) > 1$, then there is an s -ocycle for A^- .*

Proof. Small adjustments to the proof of Lemma 3.4. \square

Theorem 4.4. *Let $n, s, k \in \mathbb{Z}^+$ with $n \geq 2$, $1 \leq s \leq n-2$ and $k \leq \binom{n}{2}$.*

1. *If $\gcd(s, n) = 1$, then there is an s -ocycle for $\mathcal{W}_k(n)$.*
2. *If $\gcd(s, n) > 1$, then there is an s -ocycle for $\mathcal{W}_k^-(n)$.*

Proof. Small adjustments to the proof of Theorem 3.3. \square

Theorem 4.5. *Let $n, s, h, k \in \mathbb{Z}^+$ with $n \geq 2$, $1 \leq s \leq n-2$, $k \leq \binom{n}{2}$, and $0 \leq h < n$.*

1. *If $\gcd(s, n) = 1$, then there is an s -ocycle for $\mathcal{W}_k(n, h)$.*
2. *If $\gcd(s, n) > 1$, then there is an s -ocycle for $\mathcal{W}_k^-(n, h)$.*

Proof. Small adjustments to the proof of Theorem 3.8 \square

To produce an ocycle equivalent of Theorem 3.7, we can simplify the proof since we are now dealing with overlaps of at most $n-2$, however we require that $\gcd(s, n) = 1$.

Theorem 4.6. *For all $n, s, h \in \mathbb{Z}^+$ with $n \geq 2$, $1 \leq s \leq n-2$, $\gcd(s, n) = 1$, and $0 \leq h \leq n-1$, there is an s -ocycle for $\mathcal{W}(n, h)$.*

Proof. We define our transition graph as usual, with

$$V(G^s(n, h)) = \mathcal{W}^s(n, h) \text{ and } E(G^s(n, h)) = \mathcal{W}(n, h).$$

Clearly the graph is balanced, so we need only show that it is connected. Define the minimum vertex v^s to be the first s letters of the weak order $v = [0, h]0^{n-h-1}$. Let $x^s = x_1x_2 \dots x_s$ be an arbitrary vertex in the graph. We assume that x^s is an s -prefix of some $x = x_1x_2 \dots x_n \in \mathcal{W}(n, h)$. Applying Lemma 4.3, we may assume that x is ordered so that we have $x = [0, h]x_{h+2}x_{h+3} \dots x_n$. If $s \leq h+1$, then the first s letters of x equals v^s and we are done. Otherwise, $s \geq h+2$, and we follow the edge that corresponds to the weak order

$$[0, h]x_{h+2}x_{h+3} \dots x_s 0 \dots 0.$$

Applying Lemma 4.3 again, we can reorder this weak order to find a path to the vertex that consists of the first s letters of $[0, h]0 \dots 0x_{h+2}x_{h+3} \dots x_s$, and we are one step closer to the minimum vertex. Repeating, we will eventually arrive at the minimum vertex v^s . Thus the graph is connected, and so contains an Euler tour, and hence an s -ocycle. \square

5 Applications

Weak orders are equivalent to various other combinatorial objects. For example, they are equivalent to ordered partitions, by the following theorem

Theorem 5.1. [10], Problem 482. *There is a bijection between the set of ordered partitions of $[n]$ with the set of weak orders on $[n]$.*

Proof. An ordered partition $\mathcal{P} = \{P_0, P_1, \dots, P_k\}$ of an n -set can be defined as follows. First assume that none of the P_i 's are empty. Identify the n -set with the set $[n]$. Label the partitions by their index. Then we define a word $w = w_1w_2 \dots w_n$ to correspond to the ordered partition by setting $w_i = j$ if and only if $i \in P_j$. Using this word representation for an ordered partition, we see that every ordered partition corresponds to exactly one weak order. \square

For example, the set $\mathcal{W}(3)$ can represent the ordered partitions (corresponding to our previous listing of $\mathcal{W}(3)$ in Section 2):

$$\begin{array}{cccc} 123 & 12|3 & 1|23 & 13|2 \\ 23|1 & 2|13 & 3|12 & 1|3|2 \\ 2|1|3 & 1|2|3 & 3|1|2 & 2|3|1 \\ 3|2|1 \end{array}$$

Using this theorem, we obtain the following corollaries to our results.

Corollary 5.2. *For all $n \in \mathbb{Z}^+$, there exists a ucycle for the set of all ordered partitions of an n -set.*

Proof. This is direct from Theorem 3.2. □

Corollary 5.3. *For all $n \in \mathbb{Z}^+$ and all $h \in \mathbb{Z}^+ \cup \{0\}$ with $0 \leq h < n - 1$, there exists a universal cycle for the set of all ordered partitions of an n -set into exactly h parts.*

Proof. This is an immediate consequence of Theorem 3.7. □

Corollary 5.4. *For all n, s, h with $n \geq 2$, $1 \leq s \leq n - 2$, $\gcd(s, n) = 1$, and $h \leq n - 1$, there exists an s -ocycle for the set of all ordered partitions of an n -set into exactly h parts.*

Proof. This is immediate from Theorem 4.6. □

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